

April 2019



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Explanation of overlay networks

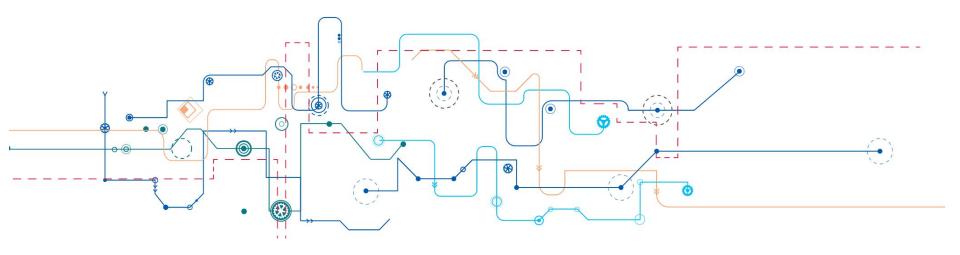
Why and how to do monitoring

Issues with monitoring overlay networks

Cubro's solution

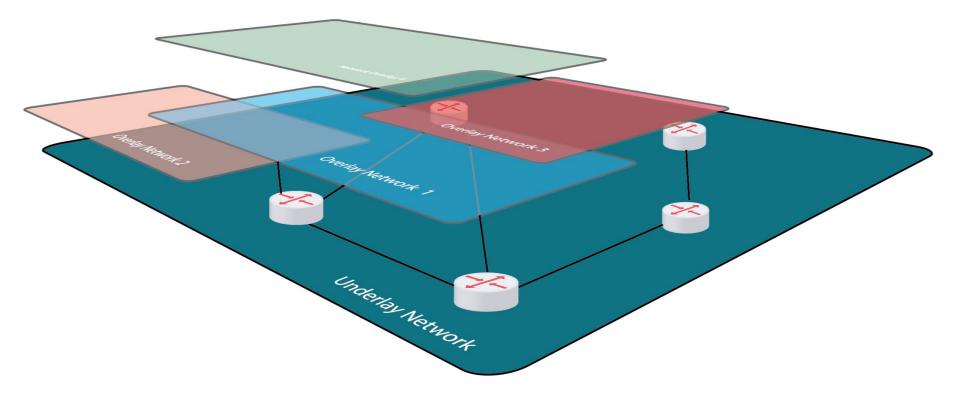
Other typically offered solutions



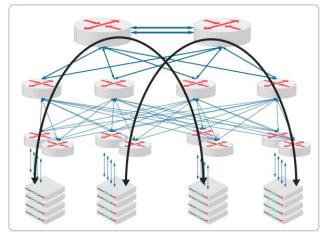


# Overlay Networks

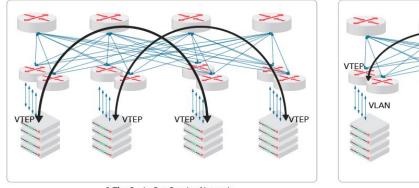
Overlay networks today are a standard in any data center, the only differentiator today versus in the past is that we are talking about hundreds or thousands of overlay networks per data center, with up to multiple thousands of endpoints. Also, overlay networks are more dynamic than in the past such that manual configuration of visibility tools is no longer possible.



# Evolution of Cloud Data Center Networking



3-Tier Overlay Network



2-Tier Scale-Out Overlay Network

EVPN - Ethernet Virtual Private Network

VTE

VTEP

VI AN

2-Tier Scale-Up Overlay Network

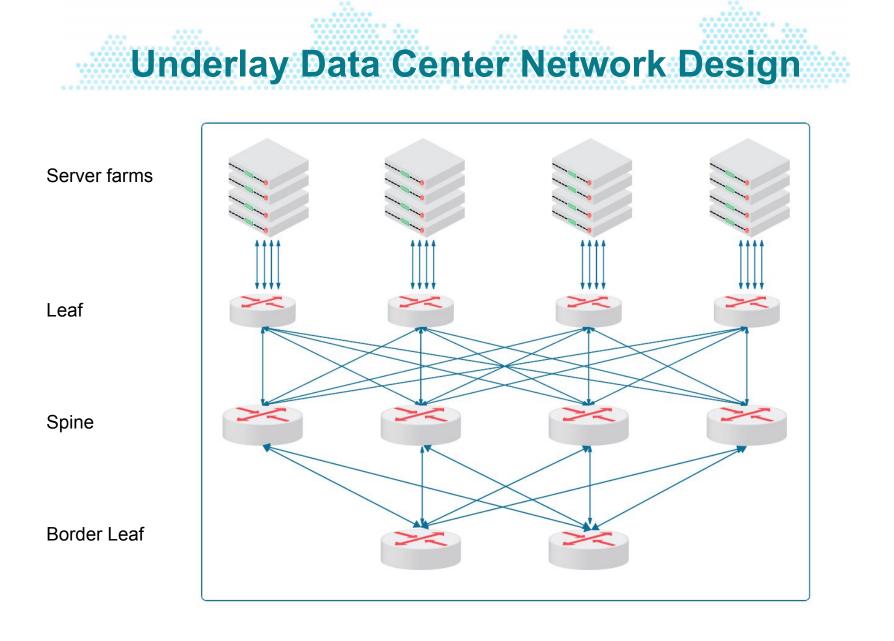
Ethernet Virtual Private Network (EVPN)[4]: The latest cloud network architecture, proposed by incumbent network mainstream players, offloads the (VXLAN) Virtual Tunnel Endpoint (VTEP) function from compute servers to leaf switches and reduces resource consumption in order to achieve a higher link throughput in compute servers.

□3-Tier Overlay Network: Same 3-tier architecture (core, aggregation and access tiers) as campus networks to provide interconnections for (VXLAN[3]) overlay traffic between compute servers.

□2-Tier Scale-Up Overlay Network: Flatten the network architecture from 3 tiers down to 2 tiers, i.e. tall spine and leaf tiers, to reduce the number of network hops (and therefore, an end to end latency) for overlay traffic between compute servers.

2-Tier Scale-Out Overlay Network: Evolved from 2-tier Scale-Up overlay network by transforming one single (usually much more expansive) tall spine switch system into multiple (more value-efficient) thin spine switch systems in order to provide:

- better redundancy: from 1+1 redundancy in the tall spine architecture to N+1 redundancy in the thin spine architecture
- better scalability: from the small cloud data centers with only a few thin spine switches to the mega cloud data centers with numerous thin spine switches
- ease of logistics: same hardware platform with the pizza-box form factor allows for inter exchange between thin spine switches and leaf switches; therefore, reduce the complexity of the cloud data center asset management and operational logistics



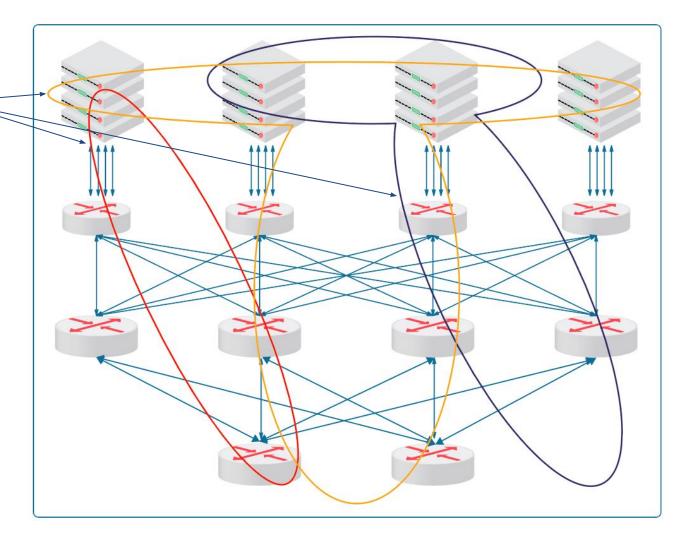
## Transparent L2 overlay network

These overlay networks share the same underlay network but for the user it is a fully transparent network.

This is good for the user of the L2 network because they can do whatever they want. For example, use any IP address or any VLAN.

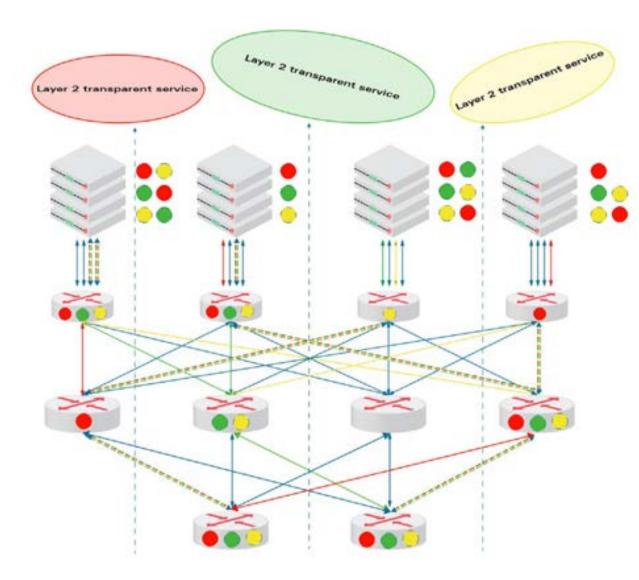
The drawback, however, is the monitoring of the underlay network because you would also see the overlay network.

Additionally, it is complex to determine the different overlay networks.



**Underlay Network Design & L2 Overlay** 





This diagram shows the general issue.

Each of these services can use the same IP ranges.

This is obviously because the guys who run these services want to make it simple for them.

The underlay network infrastructure handles the separation of these different services.

This is done by tunnels. Today this is typically VXLAN; in the old days it was MPLS or VLAN.

*The difference is today it is dynamic.* 

# Visibility Approach on Overlay & Underlay Networks

## In principle we have several visibility options

- Visibility of the underlay network.
- Visibility of a specific overlay network.
- Visibility of all overlay networks.
- Visibility of the underlay and overlay at the same time, a "full end to end view"

# Network Visibility vs Endpoint Visibility

Network Monitoring and Endpoint Monitoring are often mixed but there is a huge difference!

## Network Visibility

- Shows metrics based on network data
- Mostly passive solution
- Agnostic to devices and software
- Low operating cost
- End to end view including the transport path
- Limited application related metrics
- More complex approach in the installation phase (due to being HW)
- Good for troubleshooting

### This is Cubro's playground !

### **Endpoint Visibility**

- Show metrics based on logs or active clients
- Typically not passive
- Not agnostic, adoption for each device is usually needed.
- Shows end-to-end performance but not the network path or network parameters. The network parameters are an indirect derivation from the end-to-end parameters.
- Good application-centric metrics
- High and unpredictable operational cost
- Easy to install in the beginning
- Not efficient for troubleshooting



What is more important?

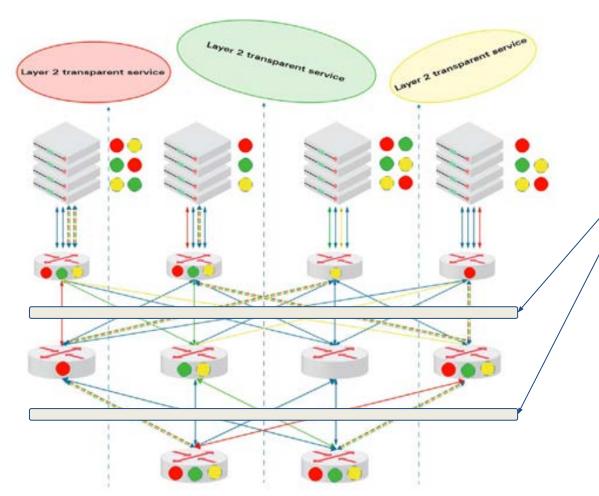
Hard to say, it depends on the customers needs... "Customers business cases"

- Service Providers = 80 % Network 20 % Endpoint
- Datacenter Providers = 90 % Network 10 % Endpoint
- Large Enterprise = 40 % Network 60 % Endpoint
- Enterprises with their own cloud infrastructure = 50 % Network 50 % Endpoint
- Enterprises with public cloud infrastructure = 10 % Network 90 % Endpoint

Bottom line: both solutions can work together to provide total visibility!

Check what you need and decide which solution works best for you!





The issue is seen clearly in this picture if you tap and monitor at these points:

We see two issues:

The same traffic can be seen several times.

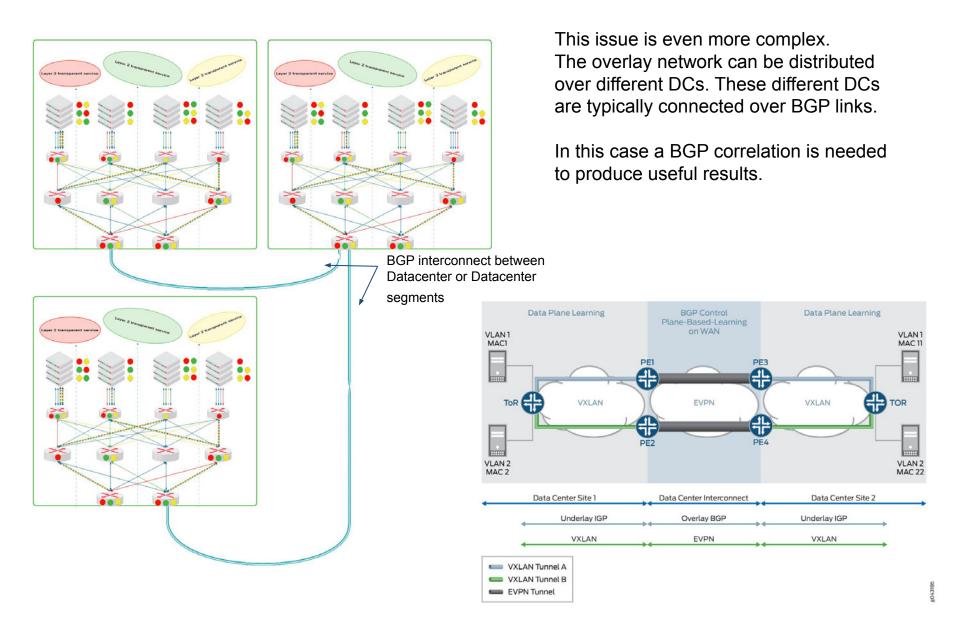
All overlay networks are seen at the same time.

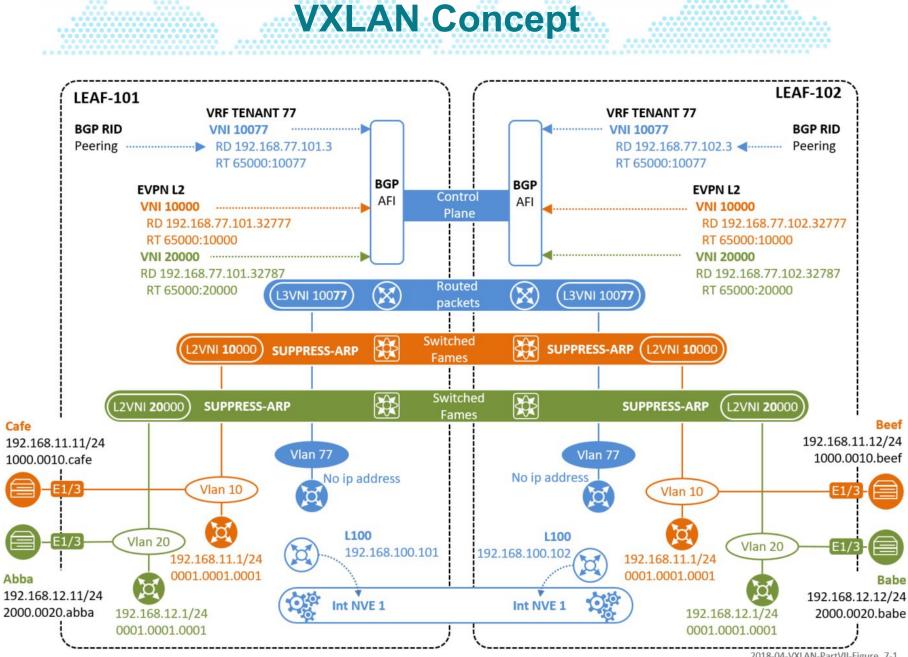
" The L2 networks can run the same IP range and therefore, it is very complex for the classical monitoring to separate the streams because typical monitoring solution works with IP addresses to determine the different paths in the network "

Typical monitoring tools cannot handle tunneled traffic.

Nearly all monitoring tools are designed to handle traffic only on one port, or on one logical network layer. This tool cannot usually correlate traffic.





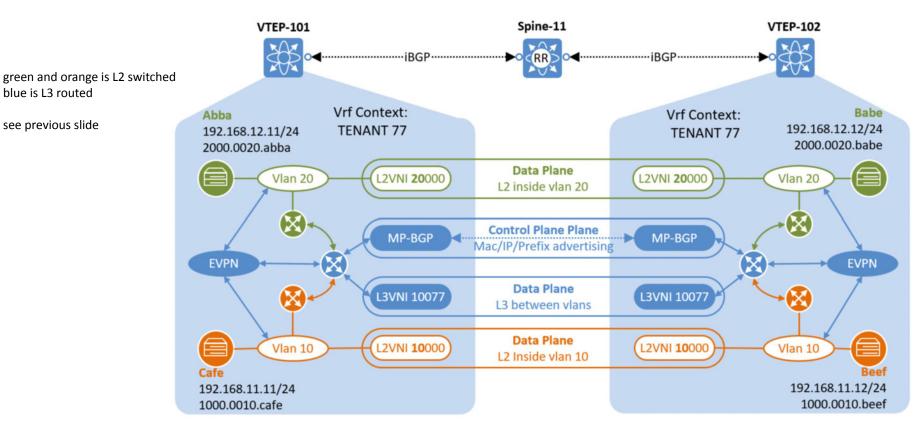


2018-04-VXLAN-PartVII-Figure 7-1



blue is L3 routed

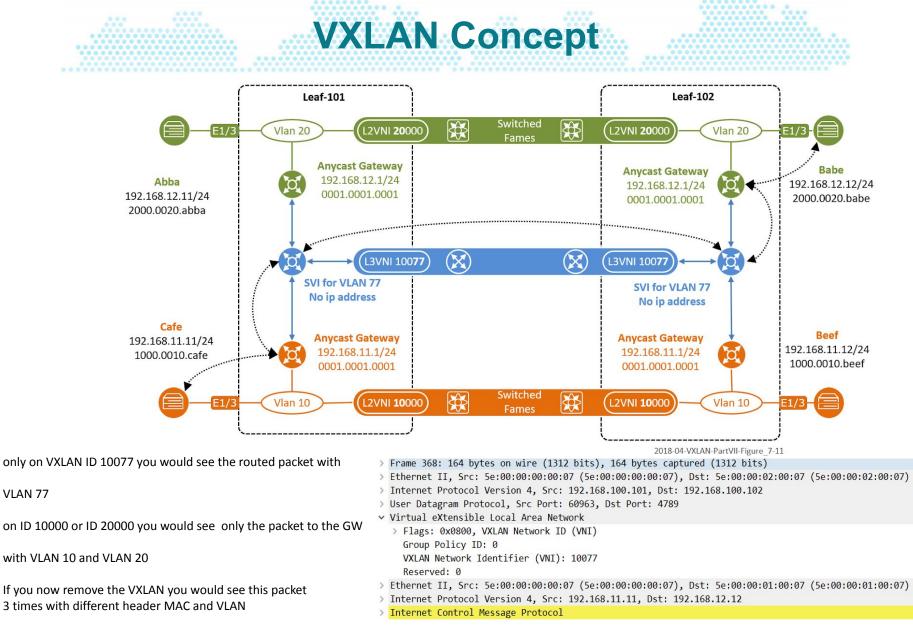
see previous slide



**VXLAN Concept** 

If you not want to monitor all traffic from Cafe, filtering on only one VXLAN is not enough, in this example you see L2 switched traffic is in VNI 10000 but routed traffic is in VNI 10077.

The challenge now is to know which VXLAN belongs together because when you have multiples routing endpoints you would have multiple VXLAN ID's. The connection is BGP !



such traffic cannot removed by a deduplication function because only the content is the same but the headers are different \*





#### **Basic connectivity test**

#### We are going to test basic connectivity between the hosts with ping.

Ping from Café to Beef (L2VNI service over VXLAN fabric)



#### Figure 7: ping Café to Beef

Cafe#pir Type esc			11.11 ce to abo	rt.				
Sending	5, 1	00-byt	ICMP Ec	hos to	192.168.11.	11, timeout	is	2 seconds:
Success	rate	is 10	percent	(5/5),	round-trip	min/avg/max	= 2	1/1/2 ms

#### Ping from Café to Abba (Local routing)



#### Figure 8: ping Café to Abba

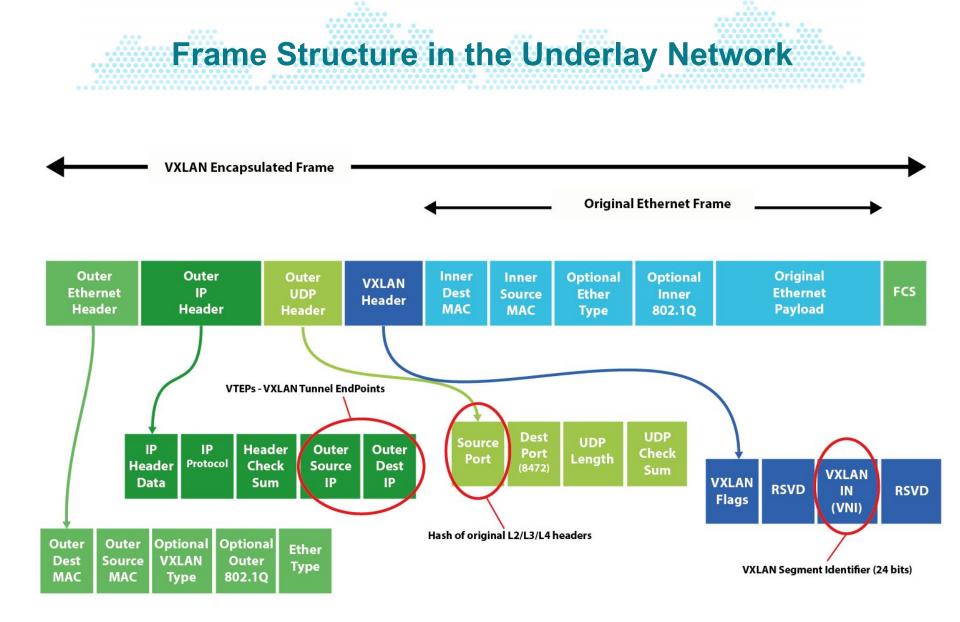
Cafe#pin Type esc Sending	ape s	equence	to			192.168.12.1	11 timeou	- 10	2	seconds.
11111	.,	0 0100	A 1/114				cay cameou	~ + ~	~	200011021
Success	rate	is 100	perce	ent (S	5/5)	, round-trip	min/avg/m	ax =	2	/8/13 ms

#### Ping from Café to Babe (L3VNI service over VXLAN fabric)



#### Figure 9: ping Café to Babe

Cafe#ping 192.168.12.12 Type escape sequence to abort. Sending 5, 100-byte ICMP Echos to 192.168.12.12, timeout is 2 seconds: !!!!! Success rate is 100 percent (5/5), round-trip min/avg/max = 20/23/29 ms



Outer	Outer	Outer	VXLAN	inner	Inner	Inner IP	Inner IP	Original	FCS
Ethernet	IP	UDP	Header	Dest	Source	Dest	Source	Ethernet	
Header	Header	Header	10020	MAC	MAC	10.10.10.10	10.10.10.1	Payload	

Issue 1

Outer	Outer	Outer	VXLAN	Inner	Inner	Inner IP	Inner IP	Original	FCS
Ethernet	IP	UDP	Header	Dest	Source	Dest	Source	Ethernet	
Header	Header	Header	10010	MAC	MAC	10.10.10.10	10.10.10.1	Payload	

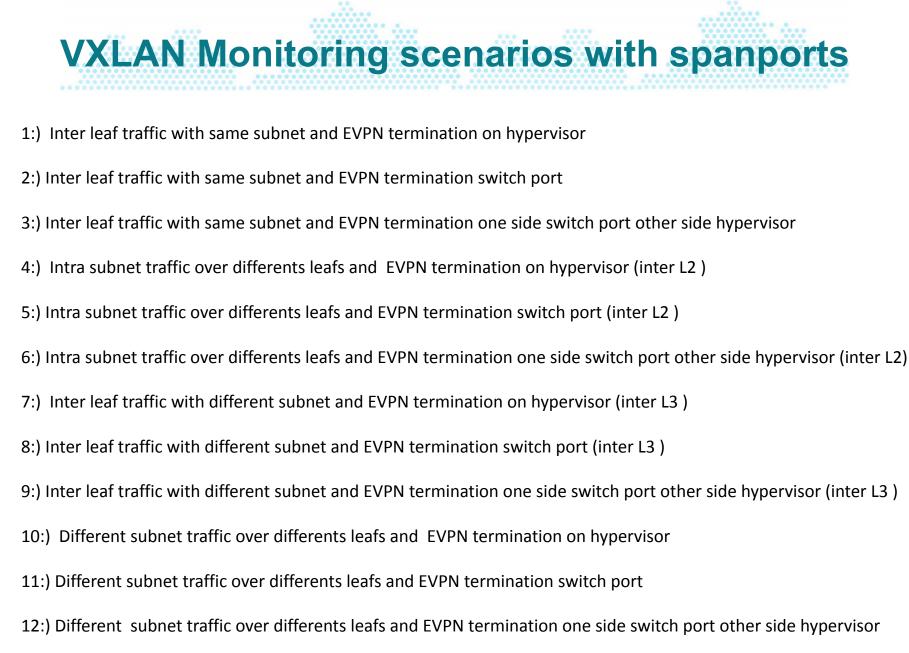
This is the common issue; same IP range but different overlay.

Normally the standard monitoring devices do not see the outer header.

Therefore, the result of the inner IP measurement is often wrong!

The overlay information is usually lost.

You get a result but it is incorrect!





1:) Inter leaf traffic with same subnet and EVPN termination on hypervisor

only egress monitoring with span port

no duplicates, but VXLAN and LAN tags on the traffic



- 2:) Inter leaf traffic with same subnet and EVPN termination switch port
  - only egress monitoring with span port
  - no duplicates, LAN tags on the traffic



- 3:) Inter leaf traffic with same subnet and EVPN termination one side switch port other side hypervisor
  - only egress monitoring with span port
  - no duplicates,
  - on the switch port termination no VXLAN but VLAN
  - on the hypervisor termination VXLAN and VLAN
  - In this case the request has no VXLAN and the answer has a VXLAN or vice versa

# VXLAN Monitoring scenarios with spanports

4:) Intra subnet traffic over differents leafs and EVPN termination on hypervisor

only egress monitoring with span port

duplicates,

packet 1:) on the spine egress VXLAN and VLAN

packet 2:) on the leaf egress VXLAN and VLAN (both the same)

remove duplicates ??

deduplication does not support VXLAN so VXLAN must be removed first very expensive approach could not be done via filters.

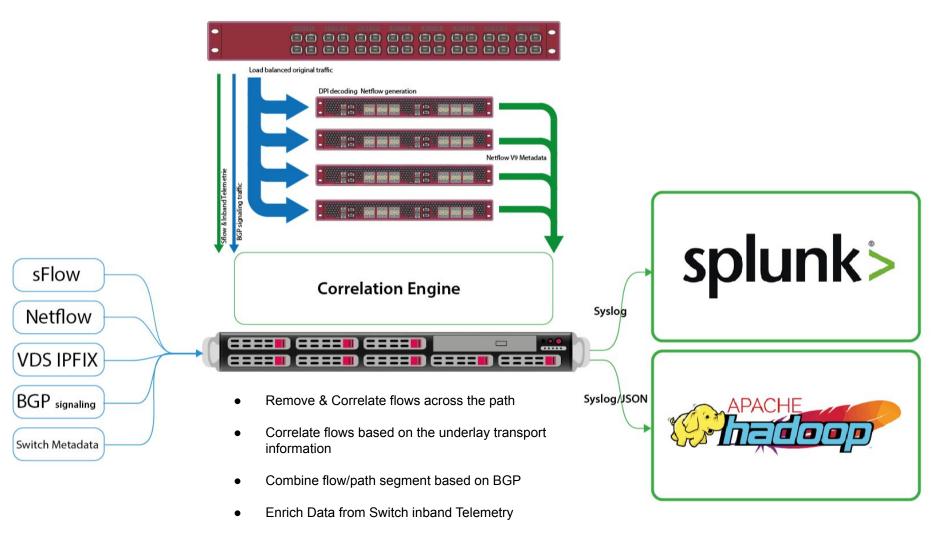
The problem is it could not be done based on links because the duplicates are on different links so first we must remove the VXLAN than aggregate all to one big pipe, but this big pipe overloads the deduplication CPU (100 Gbit +) so LB is need to feed different deduplication CPU's very complex !

(It is also not clear is the Network Layer untouched ? "IP ID field")

- 5:) Intra subnet traffic over differents leafs and EVPN termination switch port (inter L2 ) Same as 4 but no VXLAN
- 6:) Intra subnet traffic over differents leafs and EVPN termination one side switch port other side hypervisor (inter L2) Same as 4 but one packet with VXLAN + VLAN and one with only VLAN



# **Cubro Solution Design 1**



• Enrich Data with switch Table Information

### **Network Path**

Now that applications and hosts impacted by physical network outages are identified, an SDDC administrator can select end nodes to view where VM to VM traffic is encapsulated, and can see specifically which physical network devices the traffic went through.

**Cubro Solution Design 1** 



Now SDDC administrators can pinpoint which of the network devices in the path are a cause of application performance problems.

The following image shows network device interfaces involved in VM to VM communication.



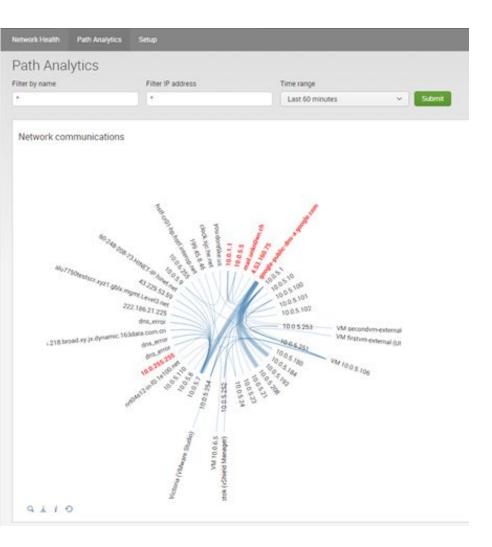
For interfaces relaying a traced communication the following information is presented:

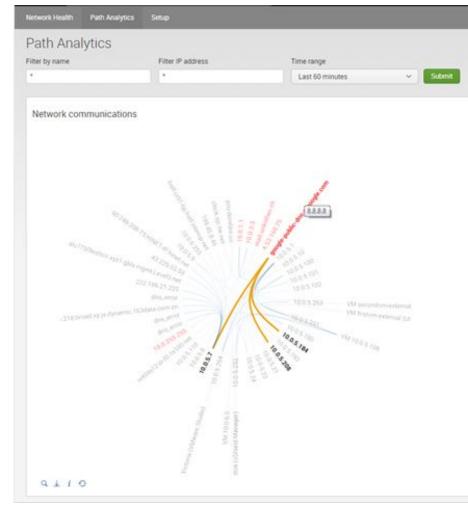
- Relative traffic load on this interface as a percent of its nominal capacity
- Relative packet rate on this interface as a percent of a maximal packet rate sustainable at a current average packet size
- A total number of bytes passed in each direction through this interface over a selected time interval
- A total number of packets passed in each direction through this interface over a selected time interval

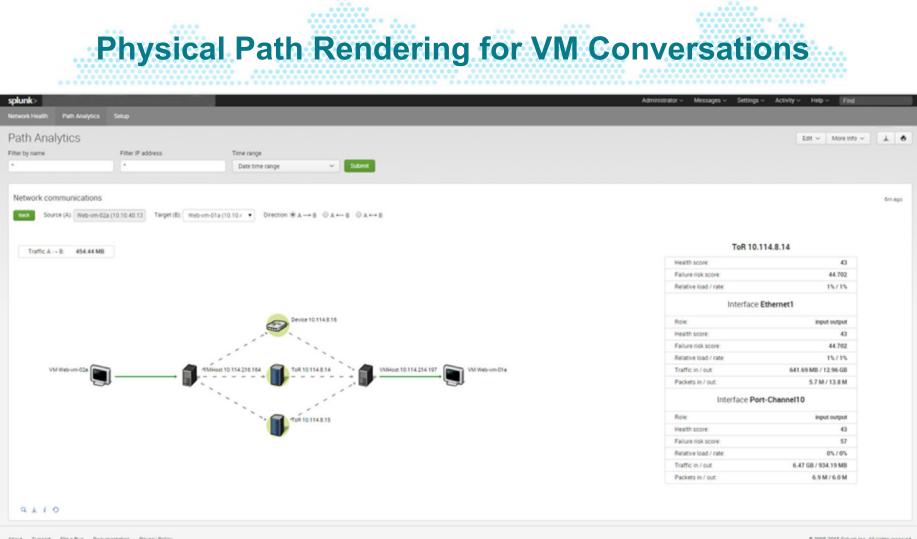
1	Health score 92 3	r
Min health score: 61	Failure risk last / worst: 5 / 39	Role: input output
Rel load / rate: 33% / 34%	Traffic in / out: 670.45 M8 / 13.38 GB	Packets in / out 6.0 M / 14.2 M
FoR 10.114.211.1	07 Interface if_2	
1	Health score 92 3	r
Min health score: 61	Failure risk last / worst: 5 / 39	Role: input output
Rel load / rate:	Traffic in / out	Packets in / out

The Path information is available not only for VM to VM (East-West traffic) within the data center, but also for VM to gateways (North-South traffic). This capability is useful in identifying network congestion and abnormal activity such as data exfiltration.

**Conversations in Virtual Overlay Networks** 



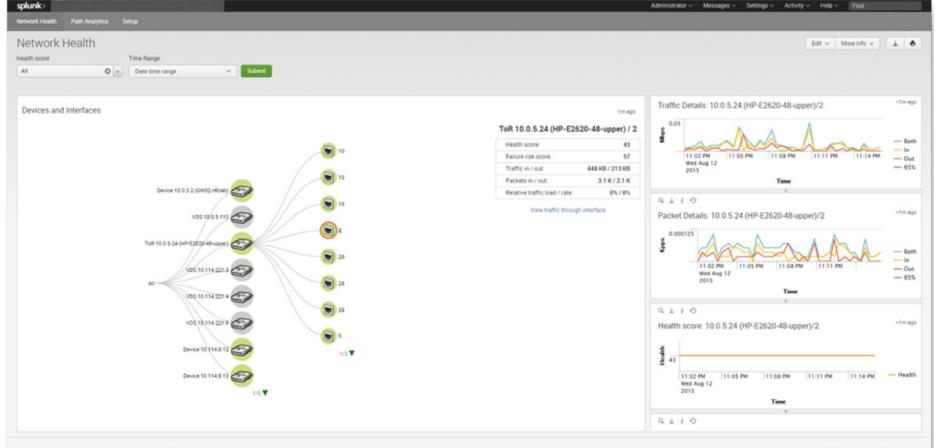




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- Top Tunnels show top tunnels by traffic
- Top Flows show top flows within a tunnel
- Distributed Firewall (DFW) coming
- Distributed Logical Router (DLR) coming

# Top Tunnels, Top Flows

• Top Tunnels (VTEPs) by Traffic; Select Time Interval; Show:

VTEP, Average Bits/s, Total Traffic Bytes, Average Packets/s, Total Packets, Total Connections

• Drill down by selecting VTEP from the list above; Show:

VXLAN\_ID, Source VM IP, Source VM Name, Source VTEP, Destination VM IP, Destination VM Name, Destination VTEP, Average Bits/s, Total Traffic Bytes, Average Packets/s, Total Packets, Total Connections

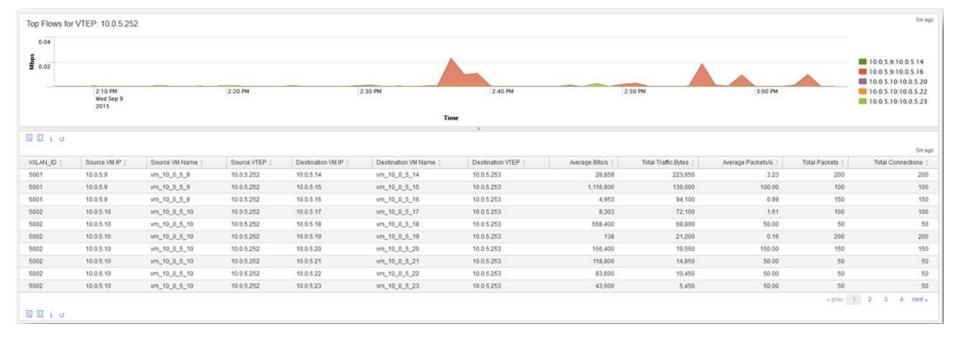


# **Top Tunnels (VTEPs)**

p Tunnels (VTEP	s)				Edit v More Info v
Range					
t 60 minutes v Sub	et .				
p Tunnels (VTEPs)					8
0.0					10.05.102 10.05.106 10.05.210 10.05.252 10.05.253 10.05.254 10.05.27
	2:10 PM 2:20 PM 2:20 PM 2:015		2.40 PM 2.50 PM	3.00 PM	10.0.5.9 199.45.8.46 98.207.85.118
		Time			0.0.5 9 199.45.8.46 96.207.85.118
	Average Bits/s	Time : Total Traffic Bytes :	Average Packets/s :	Total Packets :	10.0.5.9 199.45.8.46 96.207.85.118 Total Connection
0.5.254	Average Bits/s	Time           1         Total Traffic Bytes :           0         2,434,550	Average Packets/s : 1.24	Total Packets : 4.250	10.0.5.9 199.45.8.46 94.207.85.118 Total Connection 4
0.5.254 0.5.253	Average Bits/s	Time Total Traffic Bytes : 0 2,434,550 2 1,733,900	Average Packets/s :	Total Packets :	10.0.5.9 199.45.8.46 98.207.85.118 Total Connection 4 2
0.5.254 0.5.253 0.5.252	Average Bits/s 5.59 3.88	Time Total Traffic Bytes : 0 2.434.550 2 1.733.900 1 1.171.100	Average Packets/s : 1.24 0.76	Total Packets : 4.250 2.700	10.0.5.9 199.45.8.46 98.207.85.118 Total Connectio 4 2 2
0.5.254 0.5.253 0.5.252 0.5.7	Average Bits/s 5.59 3.88 2,76	Time Total Traffic Bytes :  2 Total Traffic Bytes :  1 1,733,900 1 1,171,100 4 1,051,400	Average Packets/s : 124 0.76 0.80	Total Packets : 4.250 2,700 2,700	10.0.5.9 199.45.8.46 98.207.85.118 Total Connectio 4 2 4 4
0.5.254 0.5.253 0.5.252 0.5.7 0.5.102	Average Bits/s 5.50 3.88 2.76 2.35	Time Total Traffic Bytes : 2 1.733,900 1.1,71,100 4 1.0,51,400 0 727,550	Average Packets/s : 124 0.76 0.80 1.27	Total Packets : 4.250 2.700 2.700 4.550	Total Connection Total Connection 4 2 2 4 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 4 2 2 2 4 2 4 2 2 2 2 2 2 2 2 2 2 2 2 2
0.5.254 0.5.253 0.5.252 0.5.7 0.5.102 0.5.106	Average Bits/s 5.90 3.88 2.76 2.35 1.76	Total Traffic Bytes:         Imme           0         2.434.550           2         1.733.900           1         1.171.100           4         1.051.400           3         727.550           4         670.400	Average Packets/s : 1.24 0.75 0.80 1.27 0.74	Total Packets : 4,250 2,700 2,700 4,550 2,450	Total Connection Total Connection Total Connection 4 2 2 2 2 2 2 2 2 2 2 2 2 2
0.5.254 0.5.253 0.5.252 0.5.7 0.5.102 0.5.105 0.5.105 9.45.8.46	Average Bits/s 5,59 3,88 2,76 2,35 1,76 1,51	Time           1         Total Traffic Bytes 1            0         2.434.550            2         1.733.900            4         1.051.400            3         727.550            4         670.400            9         633.000	Average Packets/s 1 1.24 0.76 0.80 1.27 0.74 0.62	Total Packets : 4.250 2.700 2.700 4.550 2.450 2.200	■ 100.5.9 ■ 199.45.8.46 ■ 98.207.85.118 Total Connection 4 2 2 4 2 2 2 2 2 2 2 2 2
0.5.254 0.5.253 0.5.252 0.5.7 0.5.102 0.5.105 0.5.105 0.5.210	Average Bits/s 5,59 3,88 2,76 2,35 1,76 1,51 1,42	Time           1         Total Traffic Bytes 1           0         2.434.550           2         1.733.900           1         1.717.100           4         1.051.400           3         727.550           4         670.400           9         633.000           6         594.250	Average Packets/s : 124 0.76 0.80 127 0.74 0.62 0.62	Total Packats : 4.250 2.700 4.550 2.450 2.200 2.200	10.0.5.9 199.45.8.46 94.207.85.118 Total Connection
EP : 10.5.254 10.5.253 10.5.252 10.5.102 10.5.102 10.5.102 10.5.102 10.5.252 10.5.102 10.5.252 10.5.252 10.5.252 10.5.252 10.5.252 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.254 10.5.255 10.5.25 10.5.255 1	Average Bits/s 5,50 3,88 2,76 2,35 1,76 1,51 1,42 2,82	Time           1         Total Traffic Bytes 1           0         2.434.550           2         1.733.900           1         1.717.100           4         1.051.400           3         727.550           4         670.400           9         633.000           6         594.250           2         562.300	Average Packets/s 1 124 0.76 0.80 1.27 0.74 0.62 0.62 0.62 0.62	Total Packats : 4.250 2.700 4.550 2.450 2.200 2.200 1.200	■ 100.5.9 ■ 199.45.8.46 ■ 98.207.85.118 Total Connection 4 2 2 4 2 2 2 2 1 1 1 1 1 1 1 1 1 1 1 1 1









• Top VMs by Traffic; Select Time Interval; Show:

VM\_IP, VM\_Name, Bytes\_IN, Bytes\_OUT, Packets\_IN, Packets\_OUT, Total Connections

• Drill down by selecting **Bytes\_IN** from the list above; Show (selected VM is dest\_VM):

VXLAN\_ID, src\_vm, src\_VTEP, dest\_vm, dest\_VTEP, Avg Bits/s, Bytes, Avg Packets/s, Packets, Connections

• Drill down by selecting Bytes\_OUT from the list above; Show (selected VM is src\_VM):

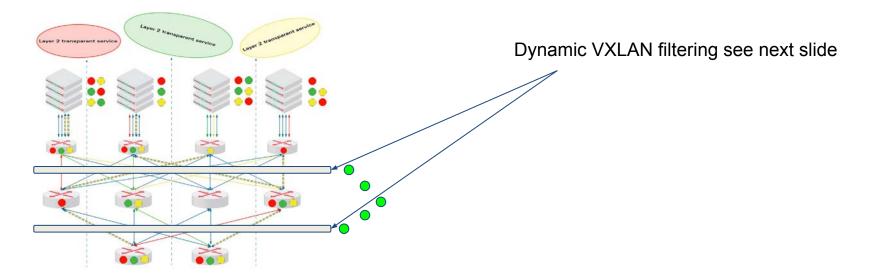
VXLAN\_ID, src\_vm, src\_VTEP, dest\_vm, dest\_VTEP, Avg Bits/s, Bytes, Avg Packets/s, Packets, Connections



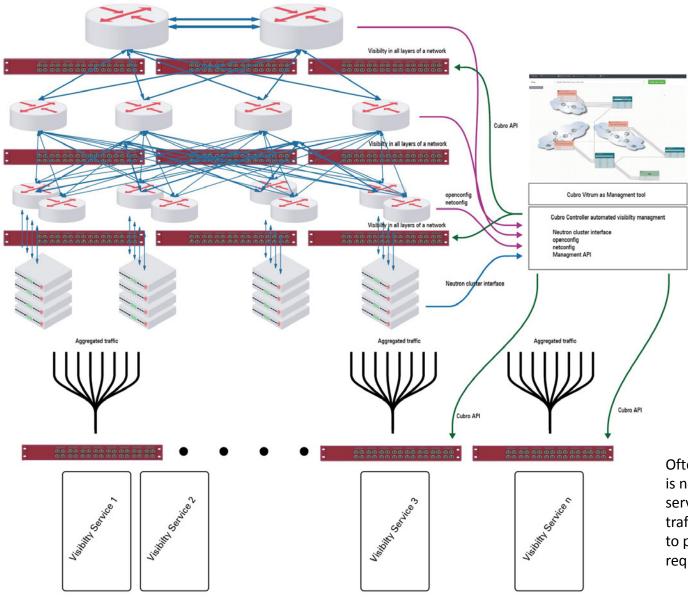
The other possible option is dynamic VXLAN filtering. This solution is needed for packet based solutions, like Wireshark and for instance mobil monitoring systems

Old equipment can be repurposed because dynamic VXLAN filtering would assure that only the traffic from the relevant overlay is filtered out and sent to legacy monitoring tools.

The challenge is that only a few NPBs are capable of VXLAN filtering. The second issue is this must be done dynamically. For that reason some signaling protocols must be decoded by the packet broker or a external appliance. -> Cubro Cloud Switch!



## Cubro Solution Design 2 the multi service approach



The center of this solution is the Cubro Controller which interacts with openstack and the other network elements, and controls the cubro visibility nodes. The cubro visibility nodes are not packet pushers they have an advanced API to cope with the visibility

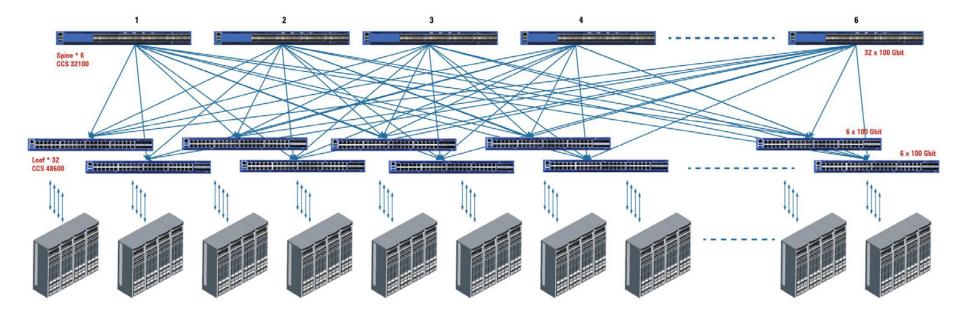
needs for today's networks!

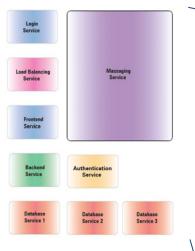
Often the same or overlapped traffic is needed for different visibility services. In this case "aggregated" traffic is send to several CSV units to provide the final preparation to the requested service.



The most advanced solution would be to use the Cubro Cloud Switch because the CCS combines an advanced switching fabric with a visibility fabric.

## In the cloud the visibility must be a part of the network infrastructure !





There is one big advantage that the application can dynamically grow when it needs more resources.

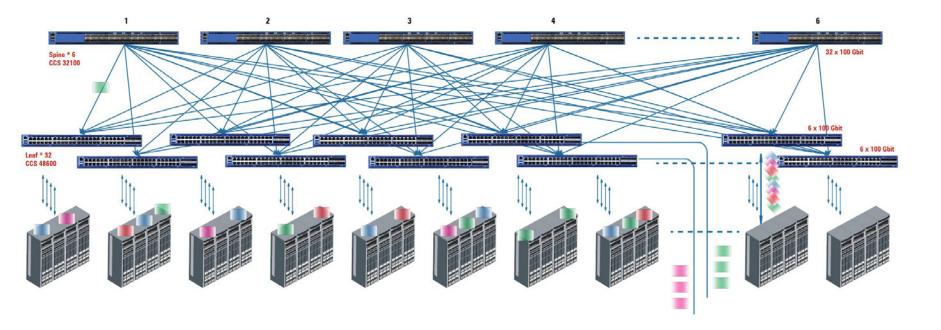
But the same "application" can also run in different data centers, spread out over the world.

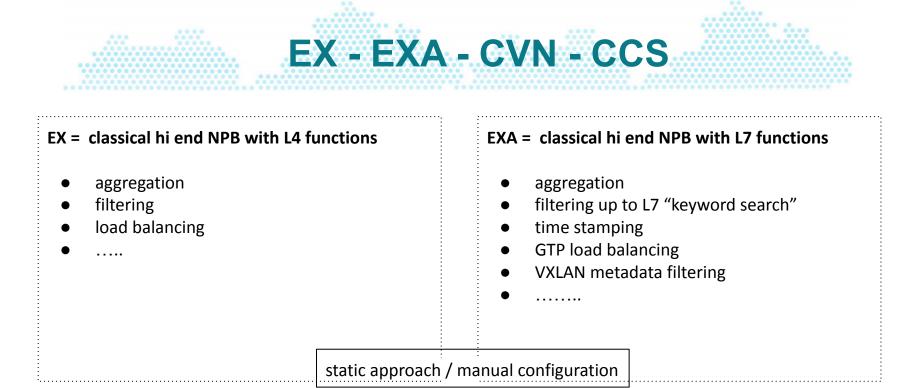
An "application" can easily breathe under load to several 1000 micro services.

**The Cloud is Breathing Application Breathing** Authenticatio Service Databa



The visibility solution had to follow the Cloud breathing in realtime and this is only possible, if the visibility solution is a part of the network infrastructure.





### CVN = Cubro Visibility Node "self organizing"

This is not a NPB any more, because it interacts with the Cubro Visibility Controller, and supports dynamic packet handling approaches for modern overlay networks.

- dynamic visibility service steering
- dynamic load balancing
- dynamic packet modification

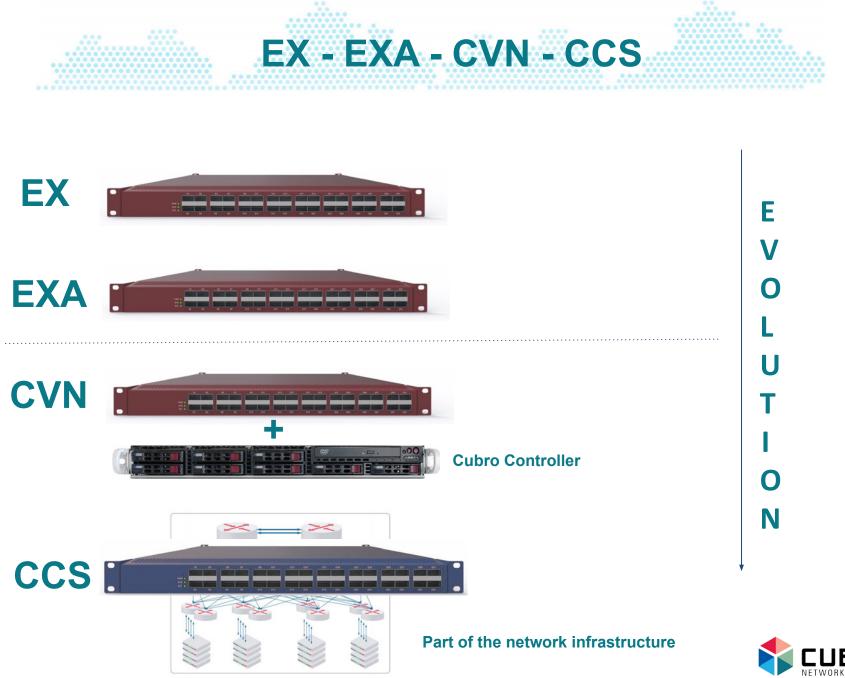
### CCS = Cubro Cloud Switch "part of the network"

The CCS is currently the end in this evolution from L4 NPB to a active network device with visibility functions included.

- inband dynamic visibility service steering
- cloud centric visibility
- application breathing support



dynamic approach / self organizing visibility









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